# Wavelets, Shearlets and Geometric Frames: Part I 

Philipp Grohs ${ }^{1}$ and Axel Obermeier²

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## Course Syllabus

MON 20th century methods (PG)
TUE 21st century methods (PG)
WED Computer class (Axel Obermeier)

## OUTLINE FOR TODAY

0. Motivation
1. Nonlinear Approximation
2. Wavelets and Point Singularities
3. Motivation

## 'Big Data'








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Source: The Economist (2010)

## The Big Challenge

Find efficient representations for complicated data

## What are Good Data Representations

Data typically modeled as function $f: \Omega \rightarrow M$, given either explicitly or (more often) implicitly as $F(f)=0$ (operator equation or inverse problem).

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If $f$ is not given explicitly, a 'good' representation should (approximately) diagonalize $F$.

## A prominent Example: JPEG 2000



Representation of image in wavelet domain

Information has its own architecture. Each data source, whether imagery, sound, text, has an inner architecture which we should attempt to discover and exploit for applications such as noise removal, signal recovery, data compression, and fast computation.
(D. Donoho, ICM Address, 2002)


## Anisotropic Data (Explicit and implicit)



Images are anisotropic!


Solutions to transport equations ( $v \cdot \nabla u+\kappa u=f$ ) are anisotropic!

## Big goal: Find optimal approximation schemes for anisotropic data!

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But What does this mean??

## 1. What's the best we can do?

A Primer in nonlinear Approximation Theory

## Start with a Mathematical Model

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For instance...
Definition (Donoho (2001))
The set $\mathcal{E}$ of cartoon images is given by

$$
\mathcal{E}:=\left\{f=f_{0}+\chi_{B} f_{1}\right\},
$$

where $f_{0}, f_{1} \in C^{2}\left([0,1]^{2}\right)$ and $\chi_{B}$ is the indicator function of $B \subset[0,1]^{2}$ with $C^{2}$ boundary.


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- Let $s^{*}(\mathcal{C}, \Phi):=\sup \left\{s>0: \mathcal{C} \subset \mathcal{A}^{s}\right\}$


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- Let $s^{*}(\mathcal{C}, \Phi):=\sup \left\{s>0: \mathcal{C} \subset \mathcal{A}^{s}\right\}$

The bigger $s^{*}(\mathcal{C}, \Phi)$, the better the dictionary $\Phi$ is suitable for the approximation of the signal class $\mathcal{C}$ !

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- © No way of efficiently computing l-term approximation of given $f$. No way of efficiently storing the dictionary $\Phi$.
$\leadsto$ need to specify the notion of 'allowed dictionaries'


## Polynomial Depth Search

Allow only 'computable' $N$-term approximations: Identify I with $\mathbb{N}$ and search best $N$-term approximation, satisfying

$$
f_{N}=\sum_{k=1}^{N} a_{\sigma(k, f)} \varphi_{\sigma(k, f)}
$$

where $\sigma(k, f) \leq \pi(k)$ for a polynomial $\pi$.

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where $\sigma(k, f) \leq \pi(k)$ for a polynomial $\pi$.
Study $s_{\text {Poly }}^{*}(\mathcal{C}, \Phi)$, defined as $s^{*}(\mathcal{C}, \Phi)$ under the constraint that the $N$-term approximations are computed using polynomial depth search and define

$$
s^{*}(\Phi):=\sup \left\{s_{\text {Poly }}^{*}(\mathcal{C}, \Phi): \Phi \text { dictionary }\right\}
$$

and use this as benchmark for the complexity of signal class $\mathcal{C}$.

## NOW WE KNOW WHAT ‘OPTIMAL’ MEANS...

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...but how can we know $s^{*}(\mathcal{C})$ ?

## Embedded Hypercubes

Definition (Donoho (2001))

1. A class of functions $\mathcal{C} \subset L^{2}\left(\mathbb{R}^{d}\right)$ contains an embedded orthogonal hypercube of dimension m and sidelength $\delta$ if there exist $f_{0} \in \mathcal{C}$ and orthogonal functions $\psi_{i} \in L^{2}\left(\mathbb{R}^{d}\right)$ for $i=1, \ldots, m$ with $\left\|\psi_{i}\right\|_{2}=\delta$ such that the collection of hypercube vertices

$$
\mathfrak{H}\left(m ; f_{0},\left(\psi_{i}\right)_{i}\right)=\left\{h=f_{0}+\sum_{i=1}^{m} \epsilon_{i} \psi_{i}: \epsilon_{i} \in\{0,1\}\right\}
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is contained in $\mathcal{C}$.
2. A class of functions $\mathcal{C}$ contains a copy of $\ell_{0}^{p}, p>0$, if there exists a sequence of orthogonal hypercubes $\left(\mathfrak{H}_{k}\right)_{k \in \mathbb{N}}$, embedded in $\mathcal{C}$, of dimensions $m_{k}$ and side-lengths $\delta_{k}$, such that $\delta_{k} \rightarrow 0$ and for $C>0$

$$
m_{k} \geq C \delta_{k}^{-p} \text { for all } k \in \mathbb{N}
$$

Theorem (Donoho (2001),
G-Keiper-Kutyniok-Schaefer (2014))
Suppose, that a class of functions $\mathcal{C} \subset L^{2}\left(\mathbb{R}^{d}\right)$ is uniformly $L^{2}$-bounded and contains a copy of $\ell_{0}^{p}$. Then

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Theorem (G-Keiper-Kutyniok-Schaefer (2014))
Let $\mathcal{C}^{\alpha}=\left\{f \in \mathcal{C}^{\alpha}[0,1]^{d}:\|f\|_{C^{\alpha}} \leq 1\right\}$. Then $\mathcal{C}$ contains $a$ copy of $\ell_{0}^{p}$ with $p=\frac{2}{2 \alpha / d+1}$. Consequently

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s^{*}\left(\mathcal{C}^{\alpha}\right) \leq \alpha / d
$$

For instance, the best we can do for $\mathcal{E}$ is an $N$-term approximation rate $N^{-1}$. We now know a benchmark result! Can we construct simple dictionaries which achieve this best-possible approximation rate??

## FRAMES

We let $\Phi$ be a frame, i.e.,

$$
A\|f\|_{2}^{2} \leq \sum_{i \in 1}\left|\left\langle f, \varphi_{i}\right\rangle\right|^{2} \leq B\|f\|_{2}^{2}
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with dual $\tilde{\Phi}=\left\{\tilde{\varphi}_{i}\right\}_{i \in 1}$.

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Definition
For $\left(c_{i}\right)_{i \in \mathbb{N}}$, define its weak $\ell^{\mathrm{P}}$-norm by

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\left\|\left(c_{i}\right)_{i \in \mathbb{N}}\right\|_{\text {wep }}:=\inf \left\{C>0:\left|c_{n}^{*}\right| \leq C n^{-1 / p}\right\},
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where $c_{n}^{*}$ is a nondecreasing rearrangement of $c_{i}$.
REMARK
We have $\left\|\left(c_{i}\right)\right\|_{\text {wep }} \leq\left\|\left(c_{i}\right)\right\|_{\ell^{\rho}}$ :

$$
\sum_{i=0}^{\infty}\left|c_{i}\right|^{p}=\sum_{i=0}^{\infty}\left|c_{i}^{*}\right|^{p} \geq \sum_{i=0}^{n}\left|c_{i}^{*}\right|^{p} \geq n\left|c_{n}^{*}\right|^{p} .
$$

## Nonlinear Approximation with Frames

Theorem
Let $\Phi$ be frame for $L^{2}$. Suppose that

$$
\left\|\left(\left\langle f, \tilde{\varphi}_{i}\right\rangle\right)_{i \in 1}\right\|_{w \ell \rho}<\infty .
$$

Then, with $s=1 / p-1 / 2$ we have

$$
f \in \mathcal{A}^{s}(\Phi)
$$

and

$$
\left\|f-\sum_{i \in l_{N}}\left\langle f, \tilde{\varphi}_{i}\right\rangle \varphi_{i}\right\|_{2} \lesssim N^{-s},
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$$

$\leadsto$ coefficient sequeces in $\ell^{p}$ for small $p$ leads to good compression! Compression works by simple coefficient thresholding!

## Literature

- DeVore. Nonlinear Approximation. Acta Numerica 7, 51-150 (1998).
- Donoho. Sparse components of images and optimal atomic decompositions. Constructive Approximation 17/3, 353-382 (2001).
- Grohs, Keiper, Kutyniok and Schaefer. Cartoon Approximation with $\alpha$-Curvelets. preprint (2014), available from www. math.ethz.ch/~pgrohs/research.
- Berger. Rate-Distortion Theory. Wiley (1971).


## 2. Wavelets and Point-Singularities

## Wavelet Approximation of Point-Singularities

Signal class

$$
\mathcal{D}^{\alpha}=\left\{f \in L^{2}[0,1]: f=f_{1} \chi_{[0, a]}+\chi_{(a, 1]} f_{2},\left\|f_{i}\right\|_{c^{\alpha}} \leq 1, a \in[0,1]\right\} .
$$

Signal


## Fourier Doesn’t Work



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## Fourier Doesn’t Work



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## Fourier Doesn’t Work

Fourier Approximation of Sawtooth Function


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Fourier Approximation of Sawtooth Function


## THEOREM

We have

$$
s^{*}\left(\mathcal{D}^{\alpha}, \text { Fourier }\right)=\frac{1}{2} .
$$

## Wavelets

Dictionaries of the form

$$
\mathcal{W}(\varphi, \psi, \alpha):=\{\underbrace{\varphi(\cdot-\alpha k)}_{\varphi_{k}}: k \in \mathbb{Z}\} \cup\{\underbrace{2^{j / 2} \psi\left(2^{j} \cdot-\alpha k\right)}_{\psi_{j, k}}: j \in \mathbb{N}, k \in \mathbb{Z}
$$

Example
Haar wavelets with $\varphi=\chi_{0,1}$ and $\psi=\frac{1}{\sqrt{2}}(\varphi(2 \cdot)-\varphi(2 \cdot-1))$. $\mathcal{W}(\varphi, \psi, 1)$ constitutes an ONB for $L^{2}$. Smooth Wavelet ONBs with compactly supported $\varphi, \psi \leadsto$ Daubechies wavelets.


Signal ( $\mathrm{N}=1024$ )



Non-Linear wavelet approximation using $\mathrm{M}=64$ coefficients


## Meyer Wavelets



Left to right: Meyer Wavelet $\psi$, Fourier Trafo of Meyer Wavelet, Fourier Trafo of Meyer Scalingfunction $\varphi$.

## Meyer Wavelets



Left to right: Meyer Wavelet $\psi$, Fourier Trafo of Meyer Wavelet, Fourier Trafo of Meyer Scalingfunction $\varphi$.

$$
\hat{\psi}_{j, k}(\xi)=2^{-j / 2} \hat{\psi}\left(2^{-j} \xi\right) \exp \left(2 \pi i 2^{-j} k \xi\right)
$$

It follows that

$$
\operatorname{supp} \psi_{j, k} \sim 2^{-j}[k-a, k+a], \quad \text { supp } \hat{\psi}_{j, k} \sim\left[2^{j-1}, 2^{j+1}\right]
$$

and thus $\left\langle f, \psi_{j, k}\right\rangle$ extracts frequency information in an annulus around $2^{j}$ and spacial information in an interval of width $2^{-j}$ around $2^{-j} k$.

## Wavelets are Optimal for Point-Singularities

## THEOREM

Suppose that $\Phi$ is the ONB of Meyer Wavelets. Then

$$
\sup _{f \in \mathcal{D}^{\alpha}}\left\|\left(\left\langle f, \psi_{j, k}\right)_{j, k}\right\rangle\right\|_{w \ell^{2 /(1+\alpha)}}<\infty
$$

In particular,

$$
s^{*}\left(\mathcal{W}(\varphi, \psi, 1), \mathcal{C}^{\alpha}\right)=s^{*}\left(\mathcal{C}^{\alpha}\right)=\alpha
$$

In other words, wavelets provide optimal approximation for functions with point-singularities, namely as efficient as if the singularity wasn't there! cproot

## Detour: Wavelets for Manifold-Valued <br> DATA <br> Many data types possess additional structure!



## Detour: Wavelets for Manifold-Valued <br> DATA

Nonlinear wavelet-type transforms can also be defined for manifold-valued data. Wavelet coefficients are elements of the tangent bundle (UrRahmen et. al., G-Wallner).

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Theorem (G 2012)
'Manifold-valued wavelets' are optimal for functions with point-singularities.

## Detour: Wavelets for Manifold-Valued DATA

Theorem (G 2012)
'Manifold-valued wavelets' are optimal for functions with point-singularities.
Proof complicated due to loss of linear structure.

## Detour: Wavelets for Manifold-Valued DATA



Left: Original entries of top 3 matrix entries of SO(3)-valued data, Right: 20:1 compression

## 2-D Wavelets

Dictionaries of the form

where

$$
\psi^{(1,1)}(x, y)=\psi(x) \psi(y), \psi^{(1,0)}(x, y)=\psi(x) \varphi(y), \psi^{(0,1)}(x, y)=\psi(y) \varphi(x), \varphi^{2 D}(x, y)=\varphi(x) \varphi(y)
$$

with $\psi$ wavelet and $\varphi$ scaling function.

## Theorem

Suppose $\mathcal{W}(\varphi, \psi, \alpha)$ is ONB or frame for $L^{2}(\mathbb{R})$. Then $\mathcal{W}^{2 D}(\varphi, \psi, \alpha)$ is ONB or frame for $L^{2}\left(\mathbb{R}^{2}\right)$.

Fourier Partitioning of 2D Meyer Wavelets


## JPEG2000

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1. Decomposing a given image $f$ in a wavelet $2 D$ wavelet basis,
2. Coefficient thesholding,
3. Quantization.

## Recall...



## Recall...



Is this really optimal??

## RECALL...

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The set $\mathcal{E}$ of cartoon images is given by

$$
\mathcal{E}:=\left\{f=f_{0}+\chi_{B} f_{1}\right\},
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where $f_{0}, f_{1} \in C^{2}\left([0,1]^{2}\right)$ and $\chi_{B}$ is the indicator function of $B \subset[0,1]^{2}$ with $C^{2}$ boundary.


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Recall benchmark approximation rate for $\mathcal{E}$ is $N^{-1}$ !

## Limitations of Wavelets

## Theorem

We have that

$$
s^{*}\left(\mathcal{E}, \mathcal{W}^{2 D}(\varphi, \psi, \alpha)\right)=\frac{1}{2} .
$$

This is one magnitude short of the optimal rate $\mathrm{N}^{-1}$.

- proof


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- proof

Can we find better dictionaries??

## Literature

- Daubechies. Ten Lectures on Wavelets. SIAM (1992)
- Mallat. A Wavelet Tour of Signal Processing. Academic Press (2008).


## End of Part I

## Appendix: Proofs and Additional Material

Proof Sketch: Suppose that $\mathfrak{h}$ is a hypercube of dimension $m$ and sidelength $\delta$.
Every element $h \in \mathfrak{h}$ can be encoded with $m$ bits.
Let $h \in \mathfrak{h}$ arbitrary and $\tilde{h}_{N}$ its best $N$-term approximation.
Assume for simplicity that $\tilde{h}_{N}$ can be encoded with $N$ bits (up to a log-factor this is possible by the polynomial depth search assumption).
Let $\hat{h}_{N}$ be the orthogonal projection of $\tilde{h}_{N}$ onto $\mathfrak{H}$.
Then we get the following coding scheme for $\{0,1\}^{m}$, i.e.
a mapping Code ${ }_{m}^{N}$ from $\{0,1\}^{m}$ to $\{0,1\}^{N}$ :

$$
\left(\epsilon_{i}\right)_{i=1}^{m} \mapsto h=f_{0}+\sum_{i=1}^{m} \epsilon_{i} \psi_{i} \in \mathfrak{h} \underbrace{\mapsto}_{N \text {-term approximation }}\left(\tau_{i}\right)_{i=1}^{N},
$$

and corresponding decoding map
$\operatorname{Decode}_{m}^{N}\{0,1\}^{N} \rightarrow\{0,1\}^{m}$

$$
\left(\tau_{i}\right)_{i=1}^{N} \underbrace{\stackrel{\leftrightarrow}{n}}_{N \text {-term reconstruction }} \tilde{h}_{N} \underbrace{\stackrel{\leftrightarrow}{n}}_{\text {projection on } \eta^{-}} \hat{h}_{N} .
$$

Result from Rate-Distortion-Theory: For $N<m / 3$ we have at least one bad $\left(\epsilon_{i}\right)_{i=1}^{m}$ (corresponding to $h \in \mathfrak{h}$ ) such that at least order $m$ bits in the reconstruction are false:
dist $_{\text {Hamming }}\left(\left(\epsilon_{i}\right)_{i=1}^{m}\right.$, Decode $\left._{m}^{N} \circ \operatorname{Encode~}_{m}^{N}\left(\epsilon_{i}\right)_{i=1}^{m}\right) \gtrsim m$.
But this implies that

$$
\left\|h-\tilde{h}_{N}\right\|_{2}^{2} \geq\left\|h-\hat{h}_{N}\right\|_{2}^{2} \gtrsim \delta^{2} m
$$

Let now $N_{k}:=m_{k} / 3$, where $m_{k}$ is the dimension of our sequence of embedded hypercubes of sidelengths

$$
\delta_{k} \gtrsim m_{k}^{-1 / p} \gtrsim N_{k}^{-1 / p}
$$

Putting this into the above equation we get

$$
\left\|h-\tilde{h}_{N_{k}}\right\|_{2}^{2} \gtrsim N_{k}^{1-2 / p} .
$$

Proof Sketch: Only $d=1$ (generalization is easy).
We need to construct a family of embedded hypercubes
( $\mathfrak{h}_{k}$ of dimensions $m_{k}$ and sidelengths $\delta_{k}$, where

$$
m_{k} \gtrsim \delta_{k}^{-p}, \quad p=\frac{2}{2 \alpha+1} .
$$

Start with $\psi \in \mathcal{C}^{\alpha}$ and

$$
\psi_{i, k}:=k^{-\alpha} \psi(k \cdot-i), \quad i, k \in \mathbb{N} .
$$

We have

$$
\left\|\psi_{i, k}\right\|_{C^{\alpha}} \leq\|\psi\|_{C^{\alpha}},
$$

so $\psi_{i, k} \in \mathcal{C}^{\alpha}$.
Furthermore, the supports of different $\psi_{i, k}, \psi_{i, k}$ are different, so for a fixed $k$, the system $\left(\psi_{i, k}\right)_{i}$ is orthogonal. A simple computation shows that

$$
\left\|\psi_{i, k}\right\|_{2}^{2}=k^{-1-2 \alpha}\|\psi\|_{2}^{2}
$$

So the system

$$
\mathfrak{h}_{k}=\left\{\sum_{i=1}^{k} \epsilon_{i} \psi_{i, k}\right\}
$$

is an embedded hypercube of dimension $m_{k}=k$ with sidelength $\delta_{k} \sim k^{-\left(\frac{1}{2}+\alpha\right)}$. ceiun to talk

Proof Sketch: Suppose for simplicity that $\Phi$ is an ONB (general case not more difficult). Let $c_{i}=\left\langle f, \varphi_{i}\right\rangle$ and assume that

$$
\left\|\left(c_{i}\right)_{i}\right\|_{w e p}<\infty .
$$

Let $\left(c_{i j}\right)_{j \in \mathbb{N}}$ be a nonincreasing rearrangement of $\left(c_{i}\right)_{i}$. The best $N$-term approximation of $f$ is given by

$$
f_{N}=\sum_{j=1}^{N} c_{j i} \varphi_{j} .
$$

We need to show that

$$
\left\|f-f_{N}\right\|_{2} \lesssim N^{-(1 / p-1 / 2)} .
$$

Since $\Phi$ is an ONB we have

$$
\left\|f-f_{N}\right\|_{2}=\left(\sum_{j=N+1}^{\infty} c_{i j}^{2}\right)^{1 / 2} .
$$

Since $\left\|\left(c_{i}\right)_{i}\right\|_{\text {wep }}<\infty$ we have that

$$
\left(\sum_{j=N+1}^{\infty} j^{-2 / p}\right)^{1 / 2}=N^{-1 / p+1 / 2}\left(N^{2 / p-1} \sum_{j=N+1}^{\infty} j^{-2 / p}\right)^{1 / 2} .
$$

So we need to show that

$$
\left(N^{2 / p-1} \sum_{j=N+1}^{\infty} j^{-2 / p}\right)^{1 / 2}<\infty
$$

Suppose for simplicity that $N=2^{\prime}$. Then

$$
\sum_{j=N+1}^{\infty} j^{-2 / p} \leq \sum_{m=1}^{\infty} F_{m}
$$

where

$$
F_{m}:=\sum_{i=2^{m-1}}^{2^{m}} i^{-2 / p} \lesssim \sum_{i=2^{m-1}}^{2^{m}} 2^{-2 m / p} \lesssim 2^{-(2 / p-1) m}
$$

So

$$
\sum_{m=1}^{\infty} F_{m} \lesssim 2^{-(2 / p-1)!}
$$

This implies the desired estimate. ©retun to talk

Proof: Let $f(x)=x, x \in[0,1]$. Then

$$
\begin{aligned}
\hat{f}(j)=\int_{0}^{1} x \exp (-2 \pi i j x) & d x=\left[x \frac{-1}{2 \pi i j} \exp (-2 \pi i j x)\right]_{0}^{1} \\
& +\int_{0}^{1} \frac{1}{2 \pi i j} \exp (-2 \pi i j x) d x=-\frac{1}{2 \pi i j}
\end{aligned}
$$

So the nonincreasing rearrangement of $c_{j}:=\langle f, \exp (2 \pi i j \cdot)\rangle$ is of the form $c_{n}^{*} \sim n^{-1}$ which lies in $w \ell^{p}$ for $p \geq 1$, which corresponds to a best $N$-term approximation order $\frac{1}{2}$.

Proof Sketch: We will show that for any $f \in \mathcal{D}^{\alpha}$ we have for every $p>\frac{1}{\alpha+1 / 2}$ that the coefficient sequence $\left\langle f, \psi_{j, k}\right\rangle$ lies in $\ell^{p}$. This is almost the result we want.
We suppose that $\frac{d^{\prime}}{d t} \hat{\psi}(0)=0$ for all $I=1, \ldots,\lfloor\alpha\rfloor$, which is definitely the case for Meyer wavelets and which is equivalent to

$$
\begin{equation*}
\int p(x) \psi(x) d x=0 \quad \text { for all polynomials } p \text { of degree } \leq\lfloor\alpha\rfloor . \tag{1}
\end{equation*}
$$

Furthermore we suppose that

$$
\begin{equation*}
\operatorname{supp} \psi_{j, l} \subset 2^{-j}[k-a, k+a] \tag{2}
\end{equation*}
$$

for some fixed $a>0$.
Now fix a scale $j$. For any bounded $f$ we have the estimate

$$
\begin{equation*}
\int f(x) \psi_{j, k}(x) d x=\int_{2^{-j}[k-a, k+a]} f(x) 2^{j / 2} \psi\left(2^{j} x-k\right) d x \lesssim 2^{-j / 2} \tag{3}
\end{equation*}
$$

For $f \in C^{\alpha}$ we have with a Taylor polynomial $p_{j, k}(x)$ of $f$ around $2^{-j} k$ that

$$
\int f(x) \psi_{j, k}(x) d x=\int_{2^{-j}[k-a, k+a]}\left(p_{j, k}(x)+O\left(2^{-\alpha j}\right)\right) 2^{j / 2} \psi\left(2^{j} x-k\right) d x .
$$

Using the vanishing-moment-property (1) and the substitution $y=2^{j} x$, we get that

$$
\begin{equation*}
\int f(x) \psi_{j, k}(x) d x \lesssim 2^{-(\alpha+1 / 2) j} \tag{4}
\end{equation*}
$$

for all $f \in C^{\alpha}$.
Let us split the index set $\mathbb{Z}$ into

$$
K_{1}=\left\{k \in \mathbb{Z}: \text { supp } \psi_{j, k} \text { intersects the singularity }\right\}
$$

and
$K_{1}=\left\{k \in \mathbb{Z}:\right.$ supp $\psi_{j, k}$ intersects the $[0,1]$ and not the singularity
By (2) we get that

$$
\begin{equation*}
\# K_{1} \lesssim 1 \text { and } \# K_{2} \lesssim 2^{j} . \tag{5}
\end{equation*}
$$

For $k \in K_{7}$ the wavelet coefficients $c_{j, k}:=\left\langle f, \psi_{j}, k\right\rangle$ satisfy $\left|c_{j, k}\right| \lesssim 2^{-j / 2}$ by (3) and for $k \in K_{2}$ we have, by (4), that $\left|c_{j, k}\right| \lesssim 2^{-j(\alpha+1 / 2)}$.
Hence we get that
$\sum_{k}\left|c_{j, k}\right|^{p} \lesssim \sum_{K_{1}} 2^{-p j / 2}+\sum_{K_{1}} 2^{-p j(\alpha+1 / 2)} \lesssim 2^{-p j / 2}+2^{-j(p \alpha+p / 2-1)}$,
by (5). Now suppose that $p>\frac{1}{\alpha+1 / 2}$. Then $\tau:=\min (p / 2, p \alpha+p / 2-1)>0$ and

$$
\sum_{j} \sum_{k}\left|c_{j, k}\right|^{D} \lesssim \sum_{j} 2^{-\tau j}<\infty
$$

Proof Sketch: Fix a scale $j$. Every wavelet $\psi_{j, k}^{\varepsilon}$ is approximately supported in a square of sidelength $\sim 2^{-j}$. We need approximately $2^{j}$ wavelets to cover the singularity curve.
For those $2^{j}$ wavelets the corresponding coefficients
$C_{j, l}^{\varepsilon}:=\left\langle f, \psi_{j, l}^{\varepsilon}\right\rangle$ satisfy

$$
\left|c_{j, l}^{\varepsilon}\right| \geq 2^{-j}
$$

So at each scale we have $\sim 2^{j}$ coefficients of magnitude $2^{-j}$.
Therefore we have

$$
\sum_{\varepsilon, k}\left|c_{j, 1}^{\varepsilon}\right|^{p} \geq \sum_{\varepsilon, k} 2^{-p j} \geq 2^{(1-p) j} .
$$

So the wavelet coefficients can only lie in $\ell^{p}$ if $p>1$, which implies best $N$-term rate of order $1 / 2$.


[^0]:    ${ }^{1}$ ETH Zürich
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